

Classification Transfer for Qualitative Reasoning Problems

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Abstract

We study formalisms for temporal and spatial reasoning in the modern, algebraic and model-theoretic, context of infinite-domain Constraint Satisfaction Problems (CSPs). We show how questions on the complexity of their subclasses can be solved using existing results via the powerful use of primitive positive (pp) interpretations and pp-homotopy. We demonstrate the methodology by giving a full complexity classification of all constraint languages that are first-order definable in Allen’s Interval Algebra and contain the basic relations (s) and (f). In the case of the Rectangle Algebra we answer in the affirmative the old open question as to whether ORD-Horn is a maximally tractable subset among the (disjunctive, binary) relations. We then generalise our results for the Rectangle Algebra to the r -dimensional Block Algebra.

1 Introduction

A classical line of enquiry within Artificial Intelligence (AI) considers an agent’s ability to reason about time and space, and a wide variety of formalisms for temporal and spatial reasoning are surveyed in [Dylla *et al.*, 2017]. Famous such examples are the Interval Algebra of [Allen, 1983], the Region Connection Calculus of [Randell *et al.*, 1992] and the Cardinal Direction Calculus of [Ligozat, 1998]. The central computational problem, designated the *fundamental task of Qualitative Spatial and Temporal Reasoning* in [Renz and Nebel, 2007], has to do with the consistency of some partial information given in one of these formalisms. In his article on Interval Algebras, Hirsch [1996] outlines *The Really Big Complexity Problem* associated to some subclass of a formalism, that is the (computational) complexity of its consistency problem. The complexity of the entire formalism is usually too blunt a measure, for example the Point Algebra [Vilain and Kautz, 1986] has long been known to be tractable, while the Interval Algebra is NP-complete [Nebel, 1995]. Much more interesting is a map of the complexity-theoretic landscape in terms of subclasses of the formalism. In the case of the Interval Algebra the natural classes of concern are subsets of disjunctions of the basic relations, and there are then *a priori* $2^{2^{13}}$ such classes requiring a complexity classification. The tower

representation produces unfortunate typography and the relevance of this number is only in its forbidding largeness, so let us replace it henceforth with a \ominus . This classification problem for the Interval Algebra tapped into a rich vein inspiring the popular paper of Nebel and Bürckert [1995], who identified a class of relations within the \ominus -many that is *maximally tractable*, in the sense that its consistency problem is in P, yet if one adds any other relation from the $\log(\ominus)$ -many the problem becomes NP-complete. This maximal class was named *ORD-Horn* and owes its tractability to a variant of the local consistency method by which Horn Satisfiability is resolved in P (a very fast algorithm for the latter is given in [Dowling and Gallier, 1984]). A complete classification showing the 18 maximal tractable classes among the \ominus -many was finally given in [Krokhin *et al.*, 2003].

A *constraint satisfaction problem* (CSP) is a computational problem in which the input consists of a finite set of variables and a finite set of constraints, and where the question is whether there exists a mapping from the variables to some fixed domain such that all the constraints are satisfied. The set of relations that is allowed to formulate the constraints in the input is called the *constraint language*. It is well-known that the consistency problem for some subclass of a temporal or spatial formalism is an example of a CSP with a constraint language over an infinite domain. Indeed, the modern and algebraic study of CSPs has played an important role in answering some of these classical questions from AI, e.g. for the Interval Algebra in [Krokhin *et al.*, 2003] (though other times, e.g. for the RCC-5, such questions were answered by exhaustive and computer-driven search through roughly \ominus -many cases [Jonsson and Drakengren, 1997]).

By now the literature on infinite-domain CSPs is beginning to mature. Much of the modern work has been motivated by questions originating in AI. The constraint languages are chosen according to what mathematical methods they might be amenable to, rather than their being idiosyncratic to spatio-temporal reasoning. These are often ordinary structures of arithmetic, but may be more elaborate mathematical constructions (such as the Fraïssé limits used for RCC-8 in [Bodirsky and Wöflf, 2011]). This need not be a hindrance, after all the Interval Algebra can be embedded in the real line and the majority of the spatio-temporal formalisms (unsurprisingly) have a natural interpretation in some Euclidean space. The notion of definability is usually in

first-order logic rather than simple disjunctions of atomic relations. Outstanding work in this direction includes the complexity classifications for the *temporal CSPs*, fo-definable in $(\mathbb{Q}; <)$ [Bodirsky and Kára, 2010] and *discrete temporal CSPs*, fo-definable in $(\mathbb{Z}; <)$ [Bodirsky *et al.*, 2015]. In these examples, a complexity-theoretic dichotomy, between P and NP-complete, is observable across the different constraint languages. Such a dichotomy is reminiscent of the case for finite-domain CSPs as detailed in the recently-proved [Bulatov, 2017; Zhuk, 2017] Feder-Vardi Conjecture [Feder and Vardi, 1999]. While infinite-domain CSPs may have much higher complexity, e.g. be undecidable [Bodirsky and Grohe, 2008], in well-behaved cases many of the methods developed for the finite-domain are applicable. Thus, the technology exists to answer many instances of the Really Big Complexity Question. Some of the tools and techniques reappear over and over again, for example tractability based on Horn’s restriction [Horn, 1951] (see [Bodirsky *et al.*, 2012]). Yet, the proofs appearing for the Interval Algebra in [Krokhin *et al.*, 2003] and the temporal CSPs of [Bodirsky and Kára, 2010] are quite distinct and often rely on local ad-hoc constructions. The observation that the Interval Algebra can be embedded in the (real or) rational line belies a new hope. Can a small number of existing classifications allow the simple derivation of other classifications? Could the classifications for the Interval Algebra and temporal CSPs be obtained, in some form, as a corollary of the other? The answer to this question is (a mildly-qualified) yes, so long as we can assume the basic relation m to be in the language defined in the Interval Algebra, and $<$ to be in the language defined in $(\mathbb{Q}; <)$. More interesting is to use existing classifications to simply derive solutions to open problems. In this vein, we will look at the Rectangle Algebra of [Guesgen, 1989; Mukerjee and Joe, 1990] and its generalisation to the r -dimensional Block Algebra of [Balbiani *et al.*, 2002].

ORD-Horn is known to be a tractable fragment not only for the Interval Algebra but also for its generalisations to the (2-dimensional) Rectangle Algebra and (r -dimensional) Block Algebra [Balbiani *et al.*, 2002]. In that paper it is noted that: “The problem of the maximality of this tractable subset [ORD-Horn] remains an open problem. Usually to prove the maximality of a fragment of a relational algebra an extensive machine-generated analysis is used. Because of the huge size ... we cannot proceed in the same way.” Our method is able to answer that ORD-Horn is indeed maximally tractable for the Rectangle Algebra as well as the r -dimensional Block Algebra, without a computational search, but via knowledge of maximal classes from [Bodirsky and Kára, 2010].

This paper is structured as follows. We introduce the preliminaries as well as the key notions of interpretations and homotopy in Section 2. Our principal results on classification transfer are given in Section 3, while we address maximal tractability of ORD-Horn in Section 4.

2 Preliminaries

Let \mathcal{B} be a finite (relational) structure, that is a domain B embellished with a finite set of relations on that domain. The sequence of arities of these relations (together with their names)

constitutes the *signature* of \mathcal{B} . Of course, one may equivalently think of \mathcal{B} as a finite set of relations over the same domain. We may now give a formal definition of the *constraint satisfaction problem* when it is parameterised by a set of relations, equivalently a (relational) structure. The problem $\text{CSP}(\mathcal{B})$, where \mathcal{B} is known as the *constraint language*, is defined as follows:

Instance: A set V of variables and a set C of constraints of the form $R(v_1, \dots, v_k)$, where k is the arity of R , $v_1, \dots, v_k \in V$ and R a relation of \mathcal{B} .

Question: Is there a function $f : V \rightarrow B$ such that $(f(v_1), \dots, f(v_k)) \in R$ for every $R(v_1, \dots, v_k) \in C$?

First-order formulas ϕ over the signature τ (or, in short, τ -formulas) are inductively defined using the logical symbols of universal and existential quantification (\forall and \exists), disjunction (\vee), conjunction (\wedge), negation (\neg), equality, bracketing, variable symbols and the symbols from τ . The semantics of a first-order formula over some τ -structure is defined in the usual Tarskian style. When ϕ is a formula without free variables (a *sentence*), we write $\mathcal{B} \models \phi$ to indicate that ϕ is true on \mathcal{B} .

One can use first-order formulas to define relations over a given structure \mathcal{B} : for a formula $\phi(x_1, \dots, x_k)$ the corresponding relation R is the set of all k -tuples $(t_1, \dots, t_k) \in B^k$ such that $\phi(t_1, \dots, t_k)$ is true in \mathcal{B} . In this case we say that R is first-order definable over \mathcal{B} . Note that our definitions are always parameter-free, i.e. we do not allow the use of domain elements in them. A structure \mathcal{A} is said to be *first-order definable* over \mathcal{B} when all of its relations are first-order definable in \mathcal{B} , it is *first-order expansion* of \mathcal{B} if it further includes all of the relations of \mathcal{B} .

The fragment of first-order logic called *primitive positive* (pp) is that which involves only existential quantification (\exists) and conjunction (\wedge). In particular, primitive positive formulas may involve equality, but no negation. An alternative definition of $\text{CSP}(\mathcal{B})$, logspace equivalent to that which we gave, is furnished by the model-checking problem for primitive positive logic, on the fixed model \mathcal{B} . In this formulation it is easy to see the following folkloric result.

Lemma 1 (Theorem 3.4 in [Jeavons, 1998]). *If a structure \mathcal{A} has a finite number of relations, all of which are pp-definable in \mathcal{B} , then there is a polynomial-time reduction from $\text{CSP}(\mathcal{A})$ to $\text{CSP}(\mathcal{B})$.*

Suppose now \mathcal{B} is a structure with a finite number of relations, each of which has arity m (this set is referred to as the *basic relations*). Define $\mathcal{B}^{\vee=}$ to contain every m -ary relation R such that $R(\bar{x})$ holds if and only if $B_1(\bar{x}) \vee \dots \vee B_t(\bar{x})$ holds, where \vee is the disjunction operator, $\{B_1, \dots, B_t\} \subseteq \mathcal{B}$, and $\bar{x} = (x_1, \dots, x_m)$ is a variable vector. Clearly, $\mathcal{B}^{\vee=}$ is a constraint language with a finite set of relations. When \mathcal{B} contains binary relations, we write $x(B_1 B_2 \dots B_t)y$ instead of $B_1(x, y) \vee B_2(x, y) \vee \dots \vee B_t(x, y)$.

2.1 The Interval and Block Algebras

The *interval algebra* [Allen, 1983] (IA) is a formalism that is both well-known and well-studied in AI. The variable domain is

$$\mathbb{I} = \{\{x \in \mathbb{Q} \mid a \leq x \leq b\} \mid a, b \in \mathbb{Q} \text{ and } a < b\},$$

Basic relation		Example	Endpoints
X precedes Y	\mathbf{p}	XXX	$X^+ < Y^-$
Y preceded by X	\mathbf{p}^\sim	YYY	
X meets Y	\mathbf{m}	XXXX	$X^+ = Y^-$
Y is met by X	\mathbf{m}^\sim	YYYY	
X overlaps Y	\mathbf{o}	XXXX YYYY	$X^- < Y^- \wedge$ $Y^- < X^+ \wedge$ $X^+ < Y^+ \wedge$
Y overlapped by X	\mathbf{o}^\sim		
X during Y	\mathbf{d}	XX	$X^- > Y^- \wedge$
Y includes X	\mathbf{d}^\sim	YYYYYY	$X^+ < Y^+$
X starts Y	\mathbf{s}	XXX	$X^- = Y^- \wedge$
Y started by X	\mathbf{s}^\sim	YYYYYY	$X^+ < Y^+$
X finishes Y	\mathbf{f}	XXX	$X^+ = Y^+ \wedge$
Y finished by X	\mathbf{f}^\sim	YYYYYY	$X^- > Y^-$
X equals Y	\equiv	XXXX YYYY	$X^- = Y^- \wedge$ $X^+ = Y^+$

Table 1: Basic relations in the interval algebra.

i.e. the variable domain consists of all closed intervals $[a, b]$ of rational numbers. If $I = [a, b] \in \mathbb{I}$, then we write I^- for a and I^+ for b . The basic relations are the 13 relations defined in Table 1. We let \mathcal{IA} denote the structure that is the set of basic interval relations over \mathbb{I} . Clearly, the 8192 relations of the IA are the contents of the set $\mathcal{IA}^{\forall=}$. Among them is the binary relation \top which holds for all pairs of intervals.

Given a sequence $S = (s_1, \dots, s_p)$, we let $S[i] = s_i$, $1 \leq i \leq p$. Let $p \geq 1$ be an integer. We will now define the p -dimensional block algebra (BA_p). The domain is \mathbb{I}^p . Given a sequence (r_1, \dots, r_p) where r_1, \dots, r_p are relations of \mathcal{IA} , we define the binary relation $B_{(r_1, \dots, r_p)} = \{(X, Y) \in (\mathbb{I}^p)^2 \mid X[i](r_i)Y[i], 1 \leq i \leq p\}$. The basic relations in BA_p constitute the structure $\mathcal{BA}_p := \{B_{(r_1, \dots, r_p)} \mid r_1, \dots, r_p \text{ from } \mathcal{IA}\}$ and the relations of BA_p are the members of $\mathcal{BA}_p^{\forall=}$. We note that BA_1 is the interval algebra and that BA_2 is often referred to as the *rectangle algebra* which we denote as RA with associated structure \mathcal{RA} . Let us note here that, for each p , $\mathcal{BA}_p^{\forall=}$ coincides with the set of binary relations first-order definable over \mathcal{BA}_p , since this latter structure admits quantifier elimination.

2.2 Interpretations

A *first-order interpretation* I of a structure \mathcal{B} over signature τ , in a structure \mathcal{A} over signature σ , is a triple (k, δ, g) where $k \in \mathbb{N}$ is the *dimension* of the reduction, δ is a first-order definable subset of \mathcal{A}^k and g is a surjection from δ to \mathcal{B} . We further require that, for every s -ary relation S in \mathcal{B} , as well as for equality, there is a first-order definable $\phi(x_1^1, \dots, x_1^k, \dots, x_s^1, \dots, x_s^k)$ over σ so that, for all $(x_1^1, \dots, x_1^k), \dots, (x_s^1, \dots, x_s^k) \in \delta$:

$$\begin{aligned} \mathcal{A} \models \phi(x_1^1, \dots, x_1^k, \dots, x_s^1, \dots, x_s^k) &\Leftrightarrow \\ \mathcal{B} \models S(g(x_1^1, \dots, x_1^k), \dots, g(x_s^1, \dots, x_s^k)). \end{aligned}$$

When ϕ is the first-order definition of R in \mathcal{B} , then S is the relation defined by ϕ_I in \mathcal{A} , where ϕ_I is built from ϕ by substituting each atomic formula in the first-order definition by

its translation under I . We will be most interested in the situation where all these fo-definitions are in fact pp-definitions, in which case we talk of a *primitive positive interpretation*.

$I_1 := (i, \delta_1, g_1), \dots, I_j := (i, \delta_j, g_j)$, of \mathcal{B} in \mathcal{A} ; together with $J := (j, \epsilon, h)$, a j -ary interpretation of \mathcal{C} in \mathcal{B} . Then we can define $J \circ (I_1, \dots, I_j)$ in the following fashion. Consider the partial ij -ary surjection $h(g_1, \dots, g_j)$, defined through $h(g_1(x_1, \dots, x_i), \dots, g_j(x_{ij-i+1}, \dots, x_{ij}))$, from \mathcal{A}^{ij} to \mathcal{C} . Note that it is defined precisely on ij -tuples satisfying $\delta_1(x_1, \dots, x_i) \wedge \dots \wedge \delta_j(x_{ij-i+1}, \dots, x_{ij}) \wedge \epsilon(g_1(x_1, \dots, x_i), \dots, g_j(x_{ij-i+1}, \dots, x_{ij}))$. Indeed, interpretations may be seen as partial surjections and partial surjections are closed under the composition we have just seen. The partial surjection $h(g_1, \dots, g_j)$ bestows a natural map from atomic relations of \mathcal{C} to relations of \mathcal{A} but there is no guarantee of nice (pp-, or even fo-) definability in \mathcal{A} . Even the definition we have given for the domain of $h(g_1, \dots, g_j)$ is not syntactic in the sense of fo-definability. However, let us now consider the case $\mathcal{C} = \mathcal{A}$ in which $h(g_1, \dots, g_j)$ is itself pp-definable, in the sense that the $ij + 1$ -ary relation on \mathcal{A} given by $h(g_1(x_1, \dots, x_i), \dots, g_j(x_{ij-i+1}, \dots, x_{ij})) = y$ is pp-definable. This is exactly $J \circ (I_1, \dots, I_j)$ being *pp-homotopic* with the identity interpretation (what we may paraphrase in future as simply “pp-homotopy”, when the context is clear)¹. In this case, it is clear that domain and relations, for $J \circ (I_1, \dots, I_j)$, are pp-definable. For example, the domain formula is $\exists y h(g_1(x_1, \dots, x_i), \dots, g_j(x_{ij-i+1}, \dots, x_{ij})) = y$ and a k -ary relation R of \mathcal{A} gives rise to its interpreting ijk -ary relation of \mathcal{A} given by

$$\begin{aligned} &\exists y^1, \dots, y^k R(y^1, \dots, y^k) \wedge \\ &h(g_1(x_1^1, \dots, x_i^1), \dots, g_j(x_{ij-i+1}^1, \dots, x_{ij}^1)) = y^1 \wedge \\ &\vdots \\ &h(g_1(x_1^k, \dots, x_i^k), \dots, g_j(x_{ij-i+1}^k, \dots, x_{ij}^k)) = y^k \end{aligned}$$

where here subscripts are read lexicographically before superscripts.

By way of example, let us consider the binary interpretation of $\mathcal{IA}^{\forall=}$ in $(\mathbb{Q}; <)$, given by $I = (2, x < y, h)$, in which $h(X^-, X^+) = [X^-, X^+]$ maps an ascending pair of rationals to the interval whose endpoints they specify. It is easy to verify that each of the relations of $\mathcal{IA}^{\forall=}$ is first-order definable in $(\mathbb{Q}; <)$. Indeed, each of the (basic) relations of \mathcal{IA} is pp-definable in $(\mathbb{Q}; <)$ (see the third column of Table 1). Thus \mathcal{IA} is even pp-interpretable in $(\mathbb{Q}; <)$.

Two structures are said to be *mutually pp-interpretable* when there is a pp-interpretation of each in the other. When these further satisfy the pp-homotopy condition then we will have enough for classification transfer.

3 Classification Transfer

Let \mathcal{C} be a structure with finite relational signature. By the *classification project* for \mathcal{C} we mean a complexity classification for $\text{CSP}(\mathcal{B})$ for all first-order expansions \mathcal{B} of \mathcal{C} that

¹We are following the terminology from [Ahlbrandt and Ziegler, 1986].

have finite relational signature. For instance, the classification project for the random graph is treated in [Bodirsky and Pinsker, 2015] and the classification project for $(\mathbb{Q}; <)$ is treated in [Bodirsky and Kára, 2010].

Sometimes, it is possible to derive the (complexity) classification project for a structure \mathcal{C} from the (complexity) classification project for another structure \mathcal{D} . The following lemma is our principal tool and gives us the method by which we can demonstrate classification transfer.

Lemma 2. *Suppose \mathcal{D} has j i -ary primitive positive interpretations I_1, \dots, I_j in \mathcal{C} , and \mathcal{C} has a j -ary primitive positive interpretation J in \mathcal{D} such that $J \circ (I_1, \dots, I_j)$ is pp-homotopic to the identity interpretation of \mathcal{C} . Then for every first-order expansion \mathcal{C}' of \mathcal{C} there is a first-order expansion \mathcal{D}' of \mathcal{D} such that \mathcal{C}' and \mathcal{D}' are mutually pp-interpretable.*

Proof. Let $I_1 = (i, U_1, g_1), \dots, I_j = (i, U_j, g_j)$ and $J = (j, V, h)$ be the primitive positive interpretations from the statement, and let \mathcal{C}' be a first-order expansion of \mathcal{C} . Then we set \mathcal{D}' to be the expansion of \mathcal{D} that contains for every k -ary atomic formula (derived from atomic relation R) in the signature of \mathcal{C}' the jk -ary relation S defined as follows. When ϕ is the first-order definition of R in \mathcal{C} , then S is the relation defined by ϕ_J in \mathcal{D} (recall that ϕ_J is built from ϕ by substituting each atomic formula in the first-order definition by its translation under J).

We claim that \mathcal{C}' is primitive positive interpretable in \mathcal{D}' , indeed by (j, V, h) . First note that V is primitive positive definable in \mathcal{D}' since \mathcal{D}' is an expansion of \mathcal{D} . An atomic formula ψ with free variables x_1, \dots, x_k in the signature of \mathcal{C}' can be interpreted in \mathcal{D}' as follows. We replace the relation symbol in ψ by its definition in \mathcal{C} , and obtain a formula ϕ in the language of \mathcal{C} . Let S be the symbol in the language of \mathcal{D}' for the relation defined by $\phi_J(x_1^1, \dots, x_1^j, \dots, x_k^1, \dots, x_k^j)$ over \mathcal{D}' . Then indeed $S(x_1^1, \dots, x_1^j, \dots, x_k^1, \dots, x_k^j)$ is a defining formula for ψ , because

$$\begin{aligned} \mathcal{C}' \models \psi(h(a_1^1, \dots, a_1^j), \dots, h(a_k^1, \dots, a_k^j)) &\Leftrightarrow \\ \mathcal{D}' \models S(a_1^1, \dots, a_1^j, \dots, a_k^1, \dots, a_k^j) & \end{aligned}$$

for all $a_1, \dots, a_k \in V$.

Conversely, we claim that \mathcal{D}' has each of the primitive positive interpretations $(i, U_1, g_1), \dots, (i, U_j, g_j)$ in \mathcal{C}' . As before, each of U_1, \dots, U_j is primitive positive definable in \mathcal{C}' since \mathcal{C}' is an expansion of \mathcal{C} . Let ϕ be a k -ary atomic formula in the (relational) signature of \mathcal{D}' . If ϕ is already in the signature of \mathcal{C} , then there is a primitive positive interpreting formula in \mathcal{C} and therefore also in \mathcal{C}' . Otherwise, by the definition of \mathcal{D}' , the relation symbol in ϕ has arity jk' , and has been introduced for a $(k \leq) k'$ -ary relation R from \mathcal{C}' (if $k < k'$ some coordinates were identified in R). We have to find a defining formula having ijk' free variables, but we will build one with ijk' variables in which some have to be identified in line with our previous comment. Let $\theta(x_0, x_1, \dots, x_{ij})$ be the primitive positive formula of arity $ij + 1$ that shows that $h(g_1(x_1, \dots, x_i), \dots, g_j(x_{ij-i+1}, \dots, x_{ij})) = x_0$ is primitive positive definable in \mathcal{C} . Then the defining formula for the atomic formula $\phi(x^1, \dots, x^{k'})$, where some of these variables are identified, is $\exists x^1, \dots, x^{k'} (R(x^1, \dots, x^{k'}) \wedge$

$\bigwedge_{\ell=1}^{k'} \theta(x^\ell, x_1^\ell, \dots, x_{ij}^\ell))$, where here subscripts are read lexicographically before superscripts. \square

Corollary 1. *Suppose \mathcal{D} has j i -ary primitive positive interpretations I_1, \dots, I_j in \mathcal{C} , and \mathcal{C} has a j -ary primitive positive interpretation J in \mathcal{D} such that $J \circ (I_1, \dots, I_j)$ is pp-homotopic to the identity interpretation of \mathcal{C} . Then for every first-order expansion \mathcal{C}' of \mathcal{C} there is a first-order expansion \mathcal{D}' of \mathcal{D} such that there are polynomial time reductions between $\text{CSP}(\mathcal{C}')$ and $\text{CSP}(\mathcal{D}')$.*

Proof. Lemma 2 tells us that for every first-order expansion \mathcal{C}' of \mathcal{C} there is a first-order expansion \mathcal{D}' of \mathcal{D} such that \mathcal{C}' and \mathcal{D}' are mutually pp-interpretable. Thus it is now incumbent on us only to argue that there are polynomial time reductions between $\text{CSP}(\mathcal{C}')$ and $\text{CSP}(\mathcal{D}')$. This essentially generalises the proof of (the hitherto unproved) Lemma 1 and is most easily given in the alternative definition of the CSP as the model-checking problem for primitive positive logic. Let $J' = (j, U, g)$ be a pp-interpretation of \mathcal{C}' in \mathcal{D}' where U is given as an j -ary pp-formula ϕ_U and each k -ary relation R of \mathcal{C}' is given as a jk -ary formula ϕ_R . From an instance of $\text{CSP}(\mathcal{C}')$, that is a pp-sentence in n variables, we build a pp-sentence in jn variables, in which each variable v becomes a sequence v^1, \dots, v^j that is constrained by ϕ_U , and where each atom of the form $R(v_1, \dots, v_k)$ is replaced by $\phi_R(v_1^1, \dots, v_1^j, \dots, v_k^1, \dots, v_k^j)$. This procedure can clearly be made in polynomial time and that it is indeed a reduction follows from it being derived from the corresponding interpretation. The argument for reducing $\text{CSP}(\mathcal{D}')$ to $\text{CSP}(\mathcal{C}')$ is dual, and the result follows. \square

3.1 The Interval Algebra

We now investigate classification transfer for the Interval Algebra.

Lemma 3. *There are 2 unary pp-interpretations, I_1 and I_2 , of $(\mathbb{Q}; <)$ in $(\mathbb{I}, \mathbf{s}, \mathbf{f})$, and a binary pp-interpretation J of $(\mathbb{I}, \mathbf{s}, \mathbf{f})$ in $(\mathbb{Q}; <)$, so that $J \circ (I_1, I_2)$ is pp-homotopic to the identity.*

Proof. Let I_1 and I_2 be $(1, \top, p_1)$ and $(1, \top, p_2)$, where p_1 and p_2 are the binary projections to start and finish point of the interval, respectively. Let J be $(2, X^- < X^+, h)$ where $h(X^-, X^+) = [X^-, X^+]$ maps pairs of points $X^- < X^+$ to the interval they specify in the line.

Define $(X < Y)_{I_1}$ and $(X < Y)_{I_2}$ by $\exists Y', W (Y(\mathbf{s})Y' \wedge X(\mathbf{s})W \wedge Y'(f)W)$ and $\exists X', W (X(f)X' \wedge Y(f)W \wedge X'(s)W)$, respectively. Define $((u_1, u_2)(\mathbf{s})(v_1, v_2))_J$ by $u_1 = v_1 \wedge u_2 < v_2$ and $((u_1, u_2)(\mathbf{f})(v_1, v_2))_J$ by $u_1 < v_1 \wedge u_2 = v_2$.

Define $(X = Y)_{I_1}$ as $X(\mathbf{s})Y$ and $(X = Y)_{I_2}$ as $X(\mathbf{f})Y$. Define $((u_1, u_2)(\equiv)(v_1, v_2))_J$ as $u_1 = v_1 \wedge u_2 = v_2$. Now $J \circ (I_1, I_2)$ maps $([X^-, X^+], [Y^-, Y^+])$ to $[Z^-, Z^+] := [X^-, Y^+]$ and pp-homotopy is given by $X(\mathbf{s})Z \wedge Y(\mathbf{f})Z$. \square

Let us make the simple observation that all the basic relations of \mathcal{IA} are pp-definable with $<$. Now we can derive the following from Corollary 1.

Corollary 2. *Let Γ be first-order definable in \mathcal{IA} containing the basic relations \mathbf{s} and \mathbf{f} . Then either $\text{CSP}(\Gamma)$ is in P or it is NP-complete.*

Noting that \mathbf{s} and \mathbf{f} are pp-definable in \mathbf{m} , though the converse is false, we note that we could have derived the previous corollary with \mathbf{m} in place of \mathbf{s} and \mathbf{f} . Corollary 2 should be compared to the main result in [Krokhin *et al.*, 2003]. They are formally incomparable, since Corollary 2 requires certain basic relations to be present, while [Krokhin *et al.*, 2003] only considers binary first-order definable relations.

3.2 The Rectangle Algebra

We now investigate classification transfer for the Rectangle Algebra. We will use the obvious fact that certain relations are pp-definable from the basic relations. Let $(\mathbf{s}, \top) := (\mathbf{s}, \mathbf{p}) \vee \dots \vee (\mathbf{s}, \equiv)$ be the disjunction with all 13 basic relations of the IA in the second coordinate. Then $X(\mathbf{s}, \top)Y$ iff $\exists Z (X(\mathbf{s}, \mathbf{p})Z \wedge Z(\mathbf{s}, \mathbf{p}^\sim)Y)$. Note also that, e.g., (\mathbf{s}, \top) is pp-definable in \mathcal{RA} by $\exists Z (X(\equiv, \mathbf{d})W \wedge Y(\equiv, \mathbf{d})W \wedge X(\mathbf{s}, \equiv)W)$.

Lemma 4. *There are 4 unary pp-interpretations, I_1, I_2, I_3 and I_4 , of $(\mathbb{Q}; <)$ in \mathcal{RA} , and a 4-ary pp-interpretation J of \mathcal{RA} in $(\mathbb{Q}; <)$, so that $J \circ (I_1, I_2, I_3, I_4)$ is pp-homotopic to the identity.*

Proof. For $i \in [4]$, let I_i be $(1, \top, p_i)$, where p_i is the 4-ary projection to the i th coordinate. Let J be $(4, X^- < X^+ \wedge Y^- < Y^+, h)$ where $h(X^-, X^+, Y^-, Y^+) = r$ maps two pairs of intervals to the rectangle they specify in the plane.

Define $(X < Y)_{I_i}$ as follows.

$$\begin{aligned} (X < Y)_{I_1} &\text{ by } \exists Y', W(Y(\mathbf{s}, \top)Y' \wedge X(\mathbf{s}, \top)W \wedge Y'(f, \top)W) \\ (X < Y)_{I_2} &\text{ by } \exists X', W(X(f, \top)X' \wedge Y(f, \top)W \wedge X'(\mathbf{s}, \top)W) \\ (X < Y)_{I_3} &\text{ by } \exists Y', W(Y(\top, \mathbf{s})Y' \wedge X(\top, \mathbf{s})W \wedge Y'(\top, f)W) \\ (X < Y)_{I_4} &\text{ by } \exists X', W(X(\top, f)X' \wedge Y(\top, f)W \wedge X'(\top, \mathbf{s})W) \end{aligned}$$

Define $((u_1, u_2, u_3, u_4)(\mathbf{s}, \top)(v_1, v_2, v_3, v_4))_J$ by $u_1 = v_1 \wedge u_2 < v_2$ and $((u_1, u_2, u_3, u_4)(f, \top)(v_1, v_2, v_3, v_4))_J$ by $u_1 < v_1 \wedge u_2 = v_2$. The other relations of \mathcal{RA} are defined in the obvious fashion.

Define $(X = Y)_{I_1}, (X = Y)_{I_2}, (X = Y)_{I_3}$ and $(X = Y)_{I_4}$ as $X(\mathbf{s}, \top)Y, X(f, \top)Y, X(\top, \mathbf{s})Y$ and $X(\top, f)Y$, respectively. Define $((u_1, u_2, u_3, u_4)(\equiv)(v_1, v_2, v_3, v_4))_J$ as $u_1 = v_1 \wedge u_2 = v_2 \wedge u_3 = v_3 \wedge u_4 = v_4$. Now, for $i \in \{1, 2, 3, 4\}$, let $W_i = (X_i, Y_i)$ where X_i and Y_i are intervals and not rectangles. $J \circ (I_1, I_2, I_3, I_4)$ maps

$$\begin{aligned} &([X_1^-, X_1^+], [Y_1^-, Y_1^+]), ([X_2^-, X_2^+], [Y_2^-, Y_2^+]), \\ &([X_3^-, X_3^+], [Y_3^-, Y_3^+]), ([X_4^-, X_4^+], [Y_4^-, Y_4^+]) \end{aligned}$$

to $Z := ([X_1^-, X_2^+], [Y_3^-, Y_4^+])$ and pp-homotopy is given by $W_1(\mathbf{s}, \top)Z \wedge W_2(f, \top)Z \wedge W_3(\top, \mathbf{s})Z \wedge W_4(\top, f)Z$. \square

Now we can derive the following from Corollary 1.

Corollary 3. *Let Γ be first-order definable in \mathcal{RA} containing the basic relations $(\mathbf{s}, \mathbf{p}), (\mathbf{s}, \mathbf{p}^\sim), (f, \mathbf{p}), (f, \mathbf{p}^\sim), (\mathbf{p}, \mathbf{s}), (\mathbf{p}^\sim, \mathbf{s}), (\mathbf{p}, f)$ and (\mathbf{p}^\sim, f) . Then either $\text{CSP}(\Gamma)$ is in P or it is NP-complete.*

4 Maximal tractability and ORD-Horn

A set of relations Γ , over the same domain, is described as *maximally tractable* among $\Delta \supseteq \Gamma$, if every finite subset of Γ gives a structure \mathcal{B} so that $\text{CSP}(\mathcal{B})$ is in P ; while for each $R \in \Delta \setminus \Gamma$, there is a finite subset of Γ given by \mathcal{B} , so that $\text{CSP}(\mathcal{B}; R)$ is NP-hard.

4.1 The Rectangle Algebra

In [Balbiani *et al.*, 2002], the strongly pre-convex relations are tied precisely to ORD-Horn, yet maximal tractability of this class, among $\mathcal{RA}^{\vee=}$, remains open. Here we settle the question by proving that ORD-Horn is indeed maximally tractable in the Rectangle Algebra. We will do this by using maximal tractability among languages first-order definable in $(\mathbb{Q}; <)$ (where in fact ORD-Horn is not maximally tractable).

A formula is called *ll-Horn* if it is a conjunction of formulas of the following form

$$\begin{aligned} (x_1 = y_1 \wedge \dots \wedge x_k = y_k) &\Rightarrow (z_1 < z_0 \vee \dots \vee z_l < z_0), \text{ or} \\ (x_1 = y_1 \wedge \dots \wedge x_k = y_k) &\Rightarrow (z_1 < z_0 \vee \dots \vee z_l < z_0 \vee \\ &(z_0 = z_1 = \dots = z_l)) \end{aligned}$$

where $0 \leq k, l$. ORD-Horn is the subclass in which we insist at most a single atom appears in each sequent. Note that k or l might be 0: if $k = 0$, we obtain a formula of the form $z_1 < z_0 \vee \dots \vee z_l < z_0$ or $(z_1 < z_0 \vee \dots \vee z_l < z_0 \vee (z_0 = z_1 = \dots = z_l))$, and if $l = 0$ we obtain a disjunction of disequalities. Also note that the variables $x_1, \dots, x_k, y_1, \dots, y_k, z_0, \dots, z_l$ need not be pairwise distinct. On the other hand, the clause $z_1 < z_2 \vee z_3 < z_4$ is an example of a formula that is *not ll-Horn*. The dual class, *dual-ll-Horn*, can be defined as *ll-Horn*, but with all $<$ replaced by $>$. The following result is from [Bodirsky and Kára, 2010] using the characterisation of [Bodirsky and Kára, 2015].

Lemma 5. *The class of relations definable in ll-Horn (resp., dual-ll-Horn) is a maximally tractable subclass of relations fo-definable in $(\mathbb{Q}; <)$.*

Lemma 6. *The relation $x = y \Rightarrow u = v$ sits in precisely two maximally tractable classes of constraint languages with respect to relations fo-definable in $(\mathbb{Q}; <)$, which are those whose relations are definable in ll-Horn and those whose relations are definable in dual-ll-Horn.*

Proof. Here we can not avoid some parlance from [Bodirsky and Kára, 2010]. Let *pp* (resp., *dual-pp*) be the binary operation on \mathbb{Q} that maps (x, y) to x , if $x < 0$, and y , otherwise (resp., maps (x, y) to y , if $x < 0$, and x , otherwise). A relation is *violated* by an operation if, when the operation is applied component-wise to some tuples in the relation, one can obtain a tuple that is not in the relation. Reading from Figure 9 in [Bodirsky and Kára, 2010], one sees that the present lemma follows if we can prove that *pp* and *dual-pp* both violate $x = y \Rightarrow u = v$. To see this for *pp*, consider the tuples $(-1, -1, 2, 2)$ and $(1, 2, 1, 2)$ for which *pp* produces the tuple $(-1, -1, 1, 2)$ which is not in the relation. The case *dual-pp* is achingly similar. \square

Call a definition in *ll-Horn* *minimal* if all of its clauses are needed and none can be replaced by one with a smaller sequent on the right-hand side of an implication. We may thus

assume that the z_0, \dots, z_l as in the definition are distinct in each clause. Note that relations definable in ORD-Horn over $(\mathbb{Q}; <)$ is a strict subset of both those relations definable in ll -Horn and those relations definable in $dual-ll$ -Horn.

Theorem 1. *The class ORD-Horn on the Rectangle Algebra is maximally tractable among the binary relations fo-definable in \mathcal{RA} .*

Proof. We know from [Balbiani *et al.*, 2002] that the ORD-Horn relations among $\mathcal{RA}^{\forall=}$ give a CSP that is tractable. Let R be some binary relation not definable in ORD-Horn and let J be as in Lemma 4. Suppose we translate R to an 8-ary relation S over $(\mathbb{Q}; <)$ via J , and let us make the similar translation for all ORD-Horn relations of the Rectangle Algebra, which will become the set of relations Γ fo-definable in $(\mathbb{Q}; <)$. Owing to Lemma 4, we need only argue that $\text{CSP}(\mathbb{Q}; \Gamma, S)$ is NP-hard (we abuse notation by presuming Γ defines also a set of relations over \mathbb{I}).

Since the (ORD-Horn) relation $\{s, \equiv, s^{\sim}\} \times (\mathbb{T} \setminus \{s, \equiv, s^{\sim}\})$ translates to $X^- = U^- \Rightarrow Y^- = V^-$ in Γ , we can deduce from Lemma 6 that the only maximally tractable classes for relations fo-definable in $(\mathbb{Q}; <)$ that $\{S\} \cup \Gamma$ can sit in are those corresponding with ll and $dual-ll$. In particular, if S is outside of these classes then it follows immediately that $\text{CSP}(\mathbb{Q}; \Gamma, S)$ is NP-hard. Let us therefore assume w.l.o.g. that S is within ll , as the other case is dual, and we will seek a contradiction.

Let ϕ be some minimal ll -Horn specification of S . Consider some clause that is not ORD-Horn. It involves some sequent of the form either $z_1 < z_0 \vee \dots \vee z_l < z_0$ or $z_1 < z_0 \vee \dots \vee z_l < z_0 \vee z_0 = z_1 = \dots = z_l$ where z_0, z_1, \dots, z_l are distinct variables. Now, since R was binary, we know that some rectangle is mentioned twice among the z_0, z_1, \dots, z_l . If some comparison involves (w.l.o.g.) I^- and I^+ , then we can remove this and we contradict minimality (note that $X^- < X^+$ because we do not allow point-like intervals). Thus, if we are not already ORD-Horn we must have something of the form, again w.l.o.g., $X^- < Y^- \vee X^- < Y^+ \vee \dots$ in the sequent, but this can be simplified to $X^- < Y^+ \vee \dots$ contradicting minimality. Thus the assumption that we be minimal actually makes us ORD-Horn. \square

4.2 The r -dimensional Block Algebra

Firstly, we will profit from studying certain automorphisms of the Block Algebra. An *automorphism* of a structure \mathcal{B} is a permutation f on its domain so that, for all relations R of \mathcal{B} , of arity k , and all k -tuples of domain elements $x_1, \dots, x_k \in B$, $R(x_1, \dots, x_k)$ iff $R(f(x_1), \dots, f(x_k))$. The Interval Algebra \mathcal{IA} enjoys all *translation* automorphisms, of the form $[X^-, X^+] \mapsto [q + X^-, q + X^+]$, for any $q \in \mathbb{Q}$. The Block Algebra \mathcal{BA}_r enjoys these similarly, independently for each of its axes. That is, \mathcal{BA}_r has all automorphisms of the form $([X_1^-, X_1^+], \dots, [X_r^-, X_r^+]) \mapsto ([q_1 + X_1^-, q_1 + X_1^+], \dots, [q_r + X_r^-, q_r + X_r^+])$. In particular, there is an automorphism of \mathcal{BA}_r that translates one axis any amount, while leaving the other axes where they are. An important property of automorphisms that we will use is that the truth of a first-order formula is invariant under an automorphism. That is, if $\phi(x_1, \dots, x_k)$ is a first-order

formula over \mathcal{BA}_r , and h is an automorphism of \mathcal{BA}_r , then $\mathcal{BA}_r \models \phi(x_1, \dots, x_k)$ iff $\mathcal{BA}_r \models \phi(h(x_1), \dots, h(x_k))$.

A *hypercuboid* is a polytope specified in k -dimensional space by the intersection of intervals $c_i \leq x_i \leq d_i$, for $c_i, d_i \in \mathbb{Q}$ and $i \in \{1, \dots, k\}$. We are now in a position to extend our classification transfers to the r -dimensional Block Algebra.

Theorem 2. *The class ORD-Horn on the r -dimensional Block Algebra is maximally tractable with respect to the binary relations fo-definable in \mathcal{BA}_r .*

Proof. We follow the proof of the Theorem 1 up to the point where we consider some clause that is not ORD-Horn. It involves some sequent of the form either $z_1 < z_0 \vee \dots \vee z_l < z_0$ or $z_1 < z_0 \vee \dots \vee z_l < z_0 \vee z_0 = z_1 = \dots = z_l$ where z_0, z_1, \dots, z_l are distinct variables. Now, since R was binary, we know that some hypercuboid is mentioned twice among the z_0, z_1, \dots, z_l .

Case A. The same dimension is mentioned twice. This has been dealt with in the proof of Theorem 1.

Case B. No dimension is mentioned twice, but we have an atom of the form $X^p < Y^q$, for $p, q \in \{+, -\}$, where I and J are different dimensions of the same hypercuboid. The r -dimensional Block Algebra enjoys an automorphism that translates the dimension associated with X while leaving unchanged all the other dimensions. Consider the evaluation that witnesses that the atom $X^p < Y^q$ is true whilst everything else in that clause is false. Now applying an automorphism we can falsify this atom while leaving all the others of that clause false (remember X^p is not repeated and $X^{\bar{p}}$, where $\bar{p} \in \{-, +\} \setminus \{p\}$, does not appear anywhere since then we would be in Case A). This demonstrates that ϕ does not specify a correct translation from the r -dimensional Block Algebra (the truth of ϕ should be invariant under these automorphisms).

Thus the assumption that we be minimal actually makes us ORD-Horn. \square

Note that our proof that ORD-Horn is maximally tractable works also for the Interval Algebra, where that result famously originated in [Nebel and Bürckert, 1995]. The only change we need to make is with Lemma 6, because $x = y \Rightarrow u = v$ can not originate from the Interval Algebra, but $x = y \Rightarrow u = v \wedge y > x \wedge v > u$ can, and suffices for our argument. Note that the analog of Corollary 3 for the Block Algebra \mathcal{BA}_r proceeds with almost identical proof.

Acknowledgements

Manuel Bodirsky has received funding from the ERC under the European Community's Seventh Framework Programme (Grant Agreement no. 681988, CSP-Infinity), and the DFG-funded project 'Homogene Strukturen, Bedingungserfüllungsprobleme, und topologische Klone' (Project number 622397). Peter Jonsson is partially supported by the *Swedish Research Council* (VR) under Grant 2017-04112. Antoine Mottet is supported by the DFG programme QuantLA.

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